

4. (*Kleinberg-Tardos 2.3*) Arrange the following list of functions in **ascending** order of growth rate. That is, if $f(n)$ precedes $g(n)$ in your list, then $f(n) = O(g(n))$. Make sure to identify any ties – $f(n)$ and $g(n)$ where $f(n) = \Theta(g(n))$.

$n^{2.5}$

$\sqrt{2n}$

$n + 10$

10^n

100^n

$n^2 \log n$

5. Imagine n students live in a dorm. You want to identify a group of k students who live in the dorm, but have never taken a class together. But, you don't know for sure that such a group exists.

- (a) How many different groups of k students are there of the n students living in the dorm?
- (b) Let's assume that k is constant (can't change). Make an argument that to loop through all such groups would be $O(n^k)$.
- (c) A brute-force algorithm to answer our question is to loop through all groups of k students and for each group, consider all pairs of students to determine if they have taken a class together. Assume that it takes constant time $O(1)$ to determine if any two students have taken a class together. What would the asymptotic runtime of this algorithm be?