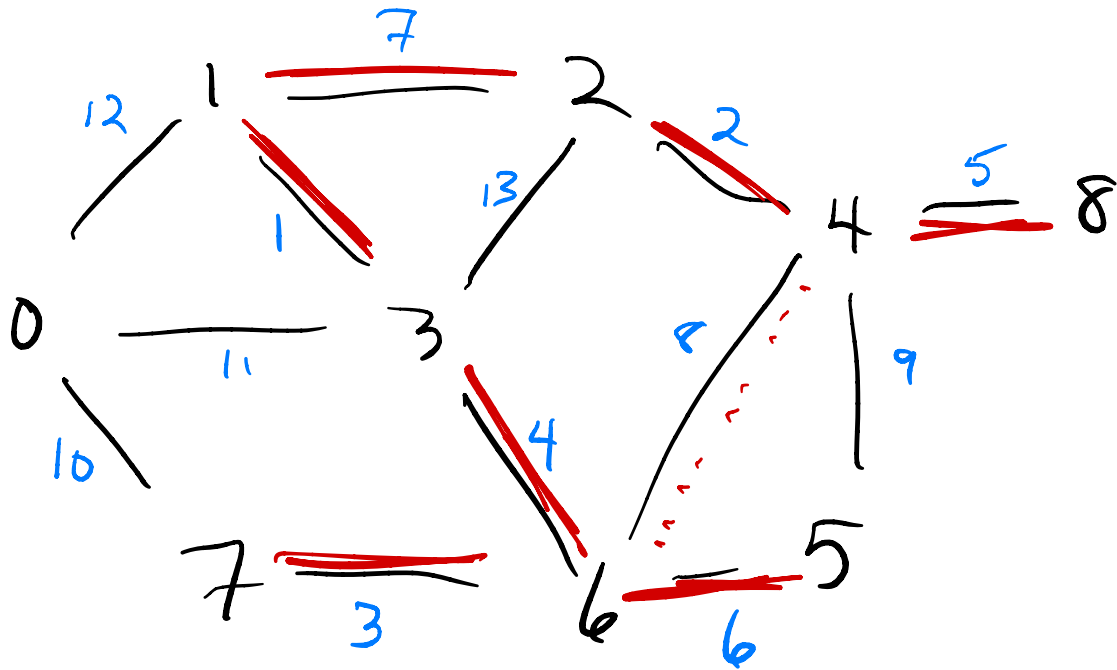
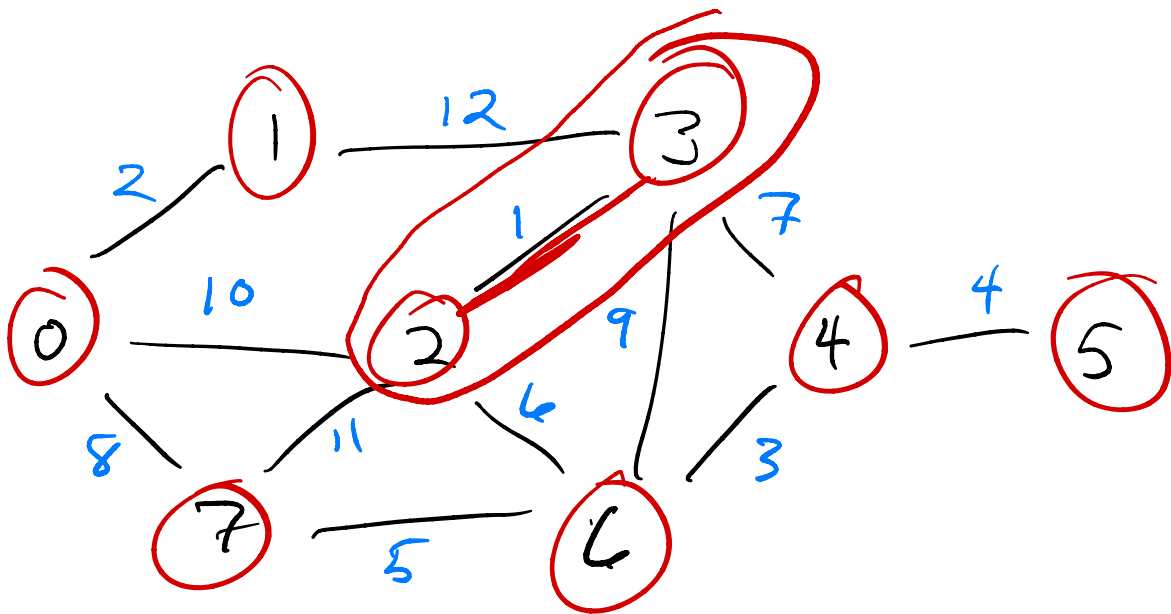


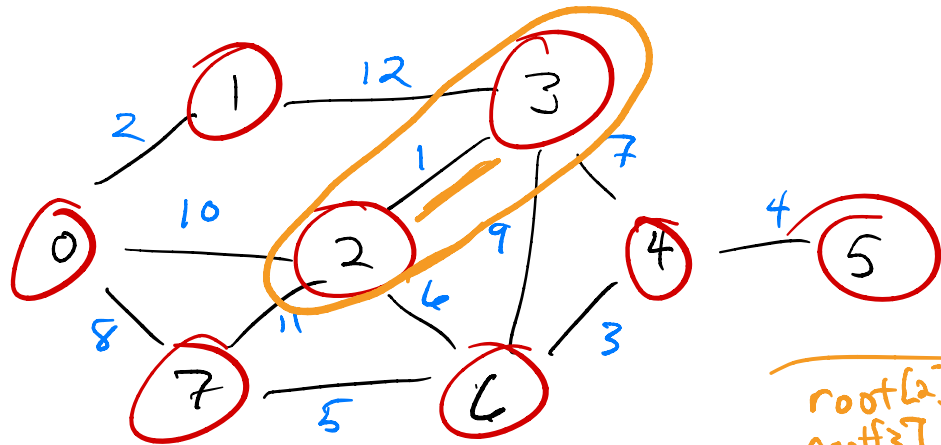
CS 252

W, 1 May 2024



Does the new edge create a cycle?
 Are the endpoints of the new edge in the same connected component?

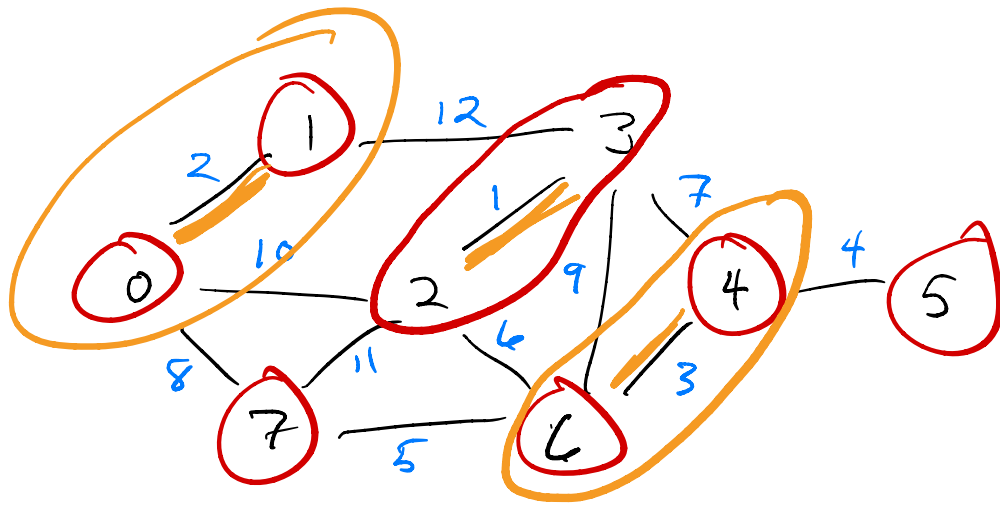




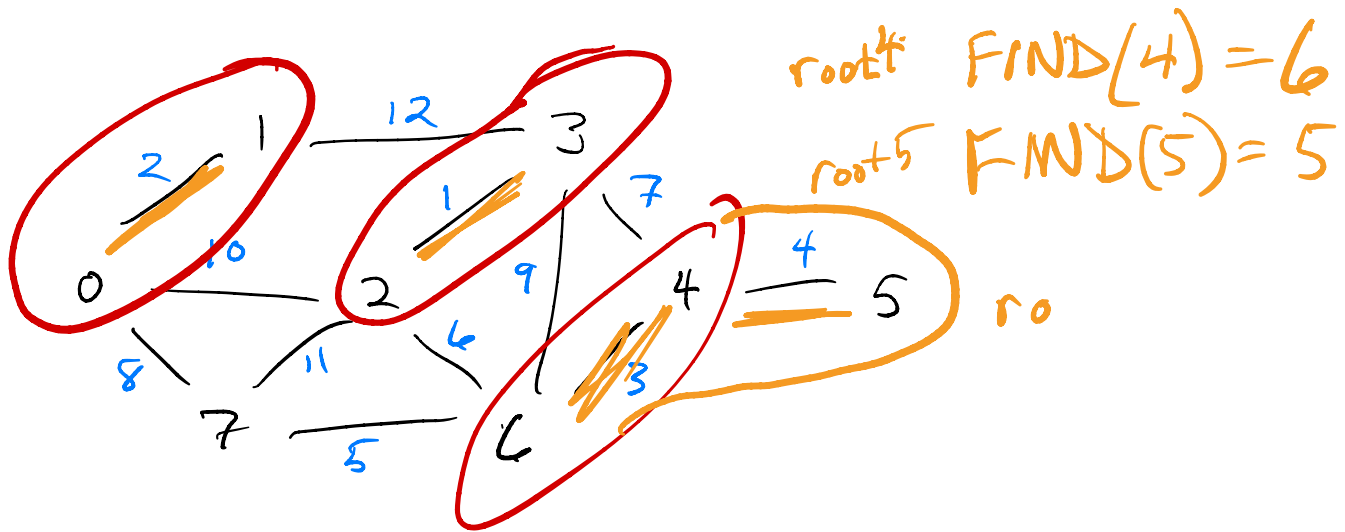
$\text{FIND}(2) == 2$
 $\text{FIND}(3) == 3$
 $\text{UNION}(2, 3)$

$\text{root}[2] = 2$ sizes
 $\text{root}[3] = 3$ both 1

	0	1	2	3	4	5	6	7
parent	0	1	2 ₃	3	4	5	6	7
size	1	1	1	1 ₂	1	1	1	1

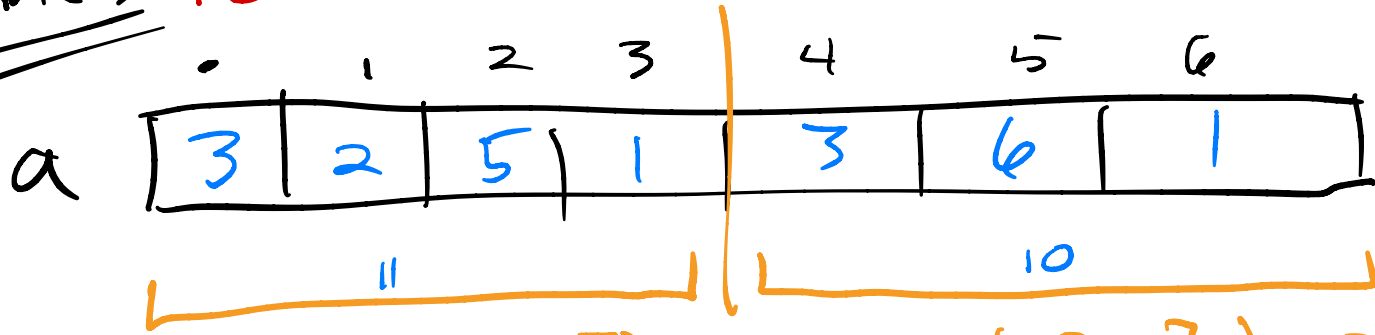


	0	1	2	3	4	5	6	7
parent	0	1	3	3	4 6	5	6	7
size	1	12	1	2	1	1	12	1



	0	1	2	3	4	5	6	7
parent	1	1	3	3	<u>6</u>	5 6	6	7
size	1	2	1	2	1	<u>1</u>	2 3	1

sum(a) $T(n)$ time



① $\text{sum}(a[0:4])$

$T(n/2)$

$\text{sum}(a[4:])$ ③

$T(n/2)$

+ ③

$O(1)$

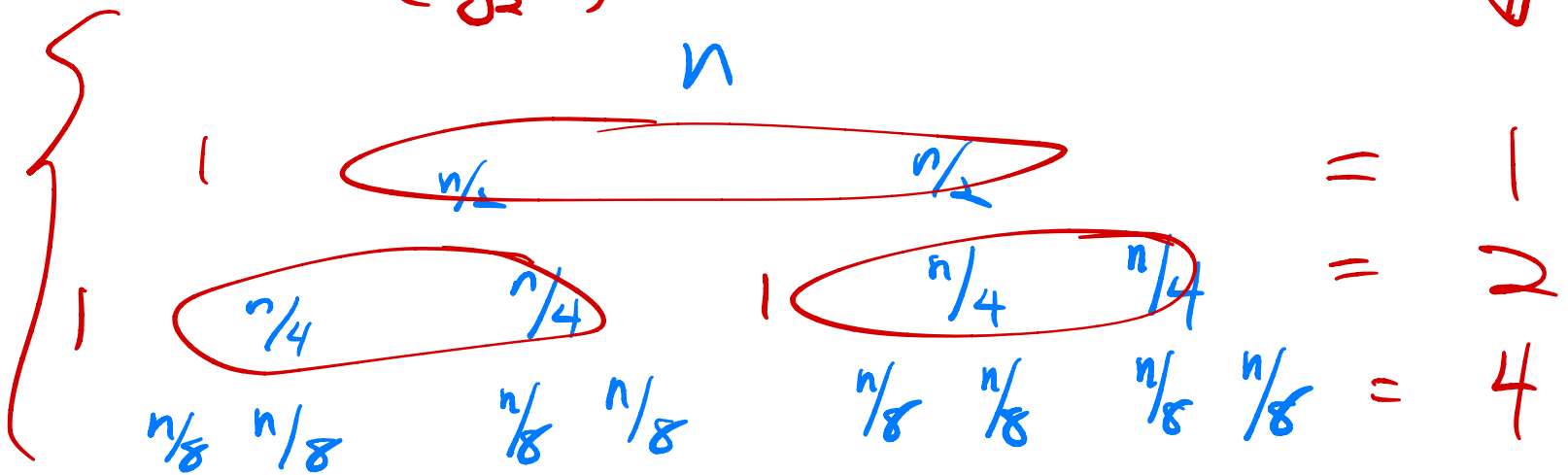
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Recursive Sum

$$T(n) \leq 2T(n/2) + 1$$

levels $\sim (\log_2 n) + 1$

Time is all here
↓



$$\begin{array}{r}
 1 \\
 + 2 \\
 + 4 \\
 + 8 \\
 \vdots \\
 + 2^{\log_2(n)-1}
 \end{array}$$

$$\begin{array}{r}
 1 + 2 + 4 + 8 \\
 \hline
 3 = 2^2 - 1 \\
 \hline
 7 = 2^3 - 1 \\
 \hline
 15 = 2^4 - 1
 \end{array}$$

$$\begin{array}{r}
 1 \\
 \hline
 n/2 \\
 \hline
 1
 \end{array}$$

$$2^{\log_2(n)-1} - 1 = \frac{2^{\log_2(n)}}{2} - 1$$

Recursive sum(a)

$$T(n) \text{ is } O(n)$$