Due Dates: One question of your choice (either #1 or #2) is due by **11:59PM Sunday, 25 September** and will be graded on completion only (1 point). Your full answers to all questions are due by **11:59PM Wednesday, 28 September**. Each solution will be graded for correctness and clarity (4 points per question). Read the course information page for further details.

At the top of your write-up for each problem, estimate the amount of time you spent on the problem, list your collaborators, and briefly describe the nature of your collaboration.

- 1. Some complexity arguments
 - (a) Define $f: \mathbb{N} \to \mathbb{R}^+$ by $f(n) = 3n^2 + 1000n + 5000$. Show that $f(n) = O(n^2)$
 - (b) Suppose $g: \mathbb{N} \to \mathbb{R}^+$ satisfies $g(n) = \Omega(n^2)$. Show that $g(n) = \Omega(n)$.
 - (c) Suppose $f, g: \mathbb{N} \to \mathbb{R}^+$ satisfy f(n) = O(n) and g(n) = O(n). Let h(m, n) = f(m) + g(n). Show that h(m, n) = O(m + n).
 - (d) Suppose $f, g: \mathbb{N} \to \mathbb{R}^+$ and $h: \mathbb{N} \times \mathbb{N} \to \mathbb{R}^+$ is defined by h(m, n) = f(m) + g(n). If h(m, n) = O(m + n) is it necessarily true that f(n) = O(n) and g(n) = O(n)? If so, prove it. If not, give a counterexample.

For the two questions involving O(m+n), you can use this definition.

Definition 1 Let $f, g : \mathbb{N} \times \mathbb{N} \to \mathbb{R}^+$. Then $\mathbf{f}(\mathbf{x}, \mathbf{y}) = \mathbf{O}(\mathbf{g}(\mathbf{x}, \mathbf{y}))$ if $\exists C > 0, N > 0$ such that $x \ge N$ and $y \ge N \implies f(x, y) \le Cg(x, y)$.

- 2. Let $S = \{s_1, s_2, ..., s_n\}$ be a set of students seeking partners for an upcoming class project. Suppose further that each s_k has created a preference list, ranking the other n-1 students. Assume that n is even.
 - (a) Provide a definition for *stable matching* in this context. (Note that among other things, your definition should handle the case where n is odd.)
 - (b) Is a given set of students and preference lists guaranteed to have a stable matching? (Here's a little hint: no.) Give a (small, I hope) example to answer this question.
 - (c) Devise an algorithm for this matching problem. It should always terminate, it should either generate a stable matching if one exists or terminate with an empty matching otherwise, and it should be as efficient as possible.
 - (d) Prove your algorithm correct and analyze its efficiency.
 - (c) Rewrite the Gale-Shapley algorithm from page 6 of Kleinberg & Tardos in a way that makes it make sense within the context of our one-set-of-students context. Change GS as little as you can.
 - (d) Show by example that your modified algorithm can result in an unstable matching.
 - (e) Show exactly where the proof of GS (i.e. the proof of claim 1.6 on pages 8-9 plus the proofs that lead up to it) fails for your modified algorithm.

Follow-up questions

As usual, I'll just put some more things to think about down here in case you have extra time, brain-space, and inclination to think about it. The grader and I won't grade anything you write about these questions, but I'm always delighted to talk with you about these and other ideas.

Anyway, here are some things you could think about. Problem 2 above concerns one of the many variations on stable matching problems. How might you adjust Gale-Shapley to deal with the following variations? What does *stable* mean in each case? How would you need to modify the algorithm? Would our original G-S proof hold up, and if not, where would it break down? Which variations correspond to what real-life applications? etc.

- There's only one set of people/items to be matched. (This is Problem 2 above.)
- The two sets to be matched have different sizes.
- *Indifference* is allowed in people's preference lists (i.e., ties are allowed—"I like Alice best, then it's a tie between Bob and Cecilia, and Doug comes last")
- People don't have to list everyone in the other set in their preference lists.
- People in M can match to multiple people in W. (This is the situation of the CS Match, where $M = \{students\}$ and $W = \{courses\}$.)
- ...?

Question 2 adapted from an assignment by Layla Oesper