

# Slithering the Link

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# Roadmap

What is a Slitherlink Puzzle?

How to Define a Slitherlink Puzzle

How to Solve a Slitherlink Puzzle

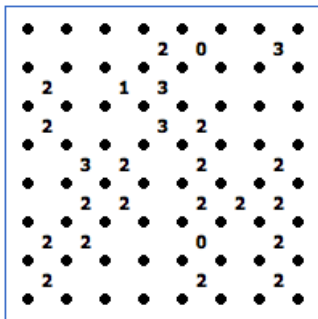
How to Make a Slitherlink Puzzle

# What is Slitherlink?

Logic puzzle developed by Nikoli

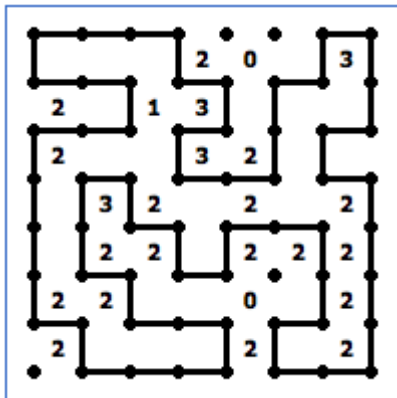
Played on:

- ▶ a rectangular lattice of dots, creating "cells"
- ▶ with some cells containing numbers



# What is Slitherlink?

Objective of the game is to create a single loop throughout the puzzle where:

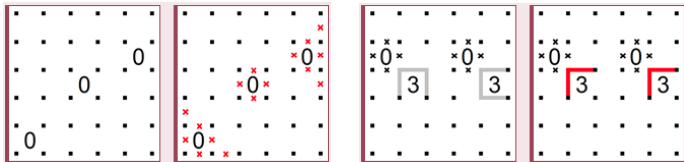


- ▶ the final solution is a continuous line that does not cross itself
- ▶ each numbered cell corresponds to the number of solution lines around it
- ▶ the puzzle should have **ONLY ONE** unique solution

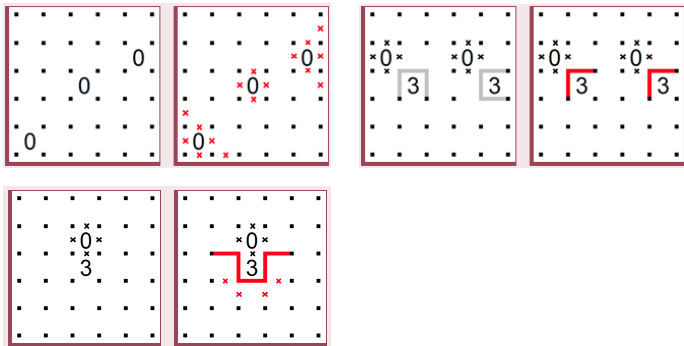




# Solving a Slitherlink Puzzle

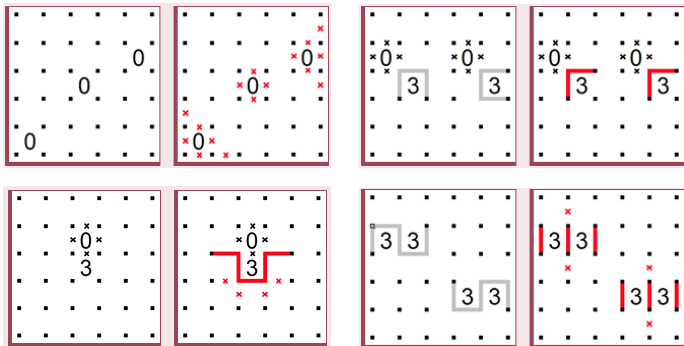


# Solving a Slitherlink Puzzle

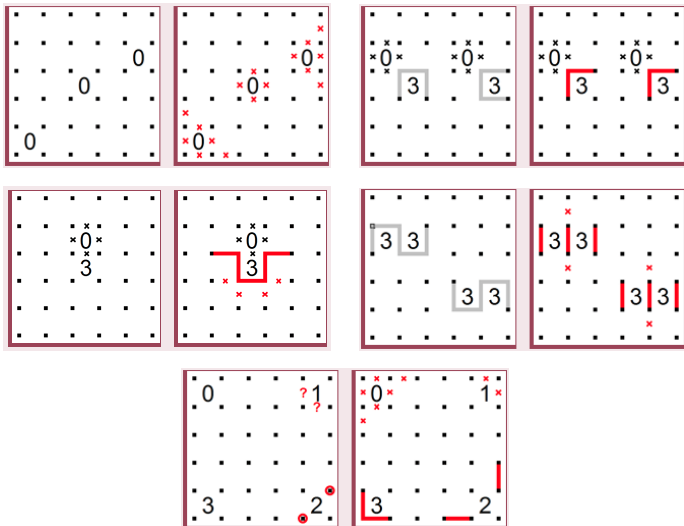




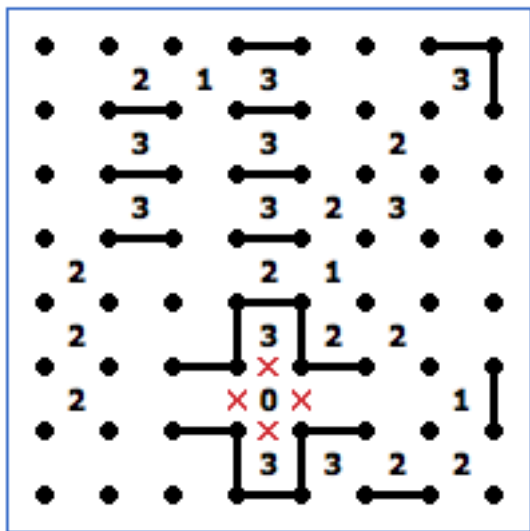
# Solving a Slitherlink Puzzle



# Solving a Slitherlink Puzzle



## Solving a Slitherlink Puzzle



## Solving a Slitherlink Puzzle

Solving a Slitherlink puzzle is an NP-complete problem, as well as determining if there are multiple solutions.

*On the NP-completeness of the Slither Link Puzzle* Takayuki YATO

# Puzzle Representation

Components of a puzzle:

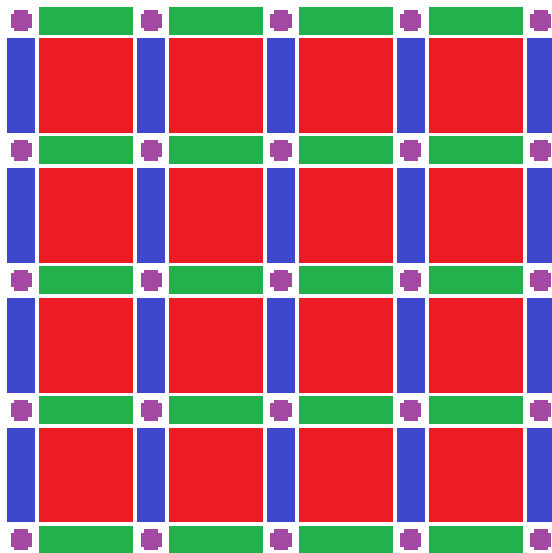
- ▶ the grid
  - ▶ lines
  - ▶ numbers
- ▶ rules and contradictions
- ▶ contours
- ▶ what it means to be solved

# The Grid

M by N grid has 3 2D arrays:

- ▶ M by N 2D array for numbers; values 0 to 3 or empty
- ▶ M + 1 by N 2D array for horizontal lines; values line, x, or empty
- ▶ M by N + 1 2D array for vertical lines; values line, x, or empty

*A rule-based approach to the puzzle of Slither Link.* Stefan Herting.



# Rules and Contradictions

Each rule has:

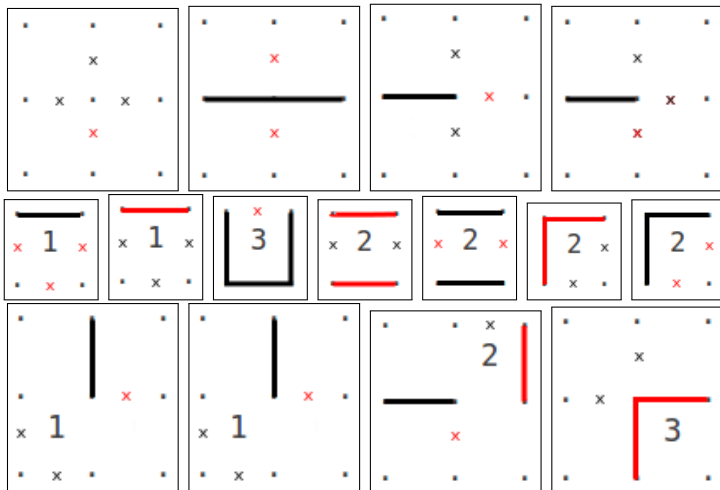
- ▶ dimensions
- ▶ prerequisites
- ▶ consequences

Each contradiction has:

- ▶ dimensions
- ▶ prerequisites

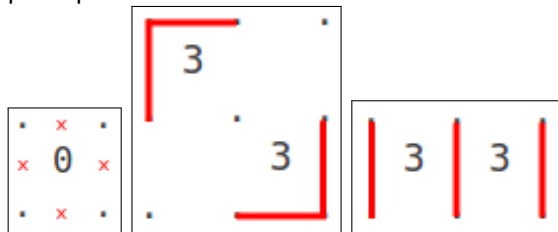


# Examples of Rules

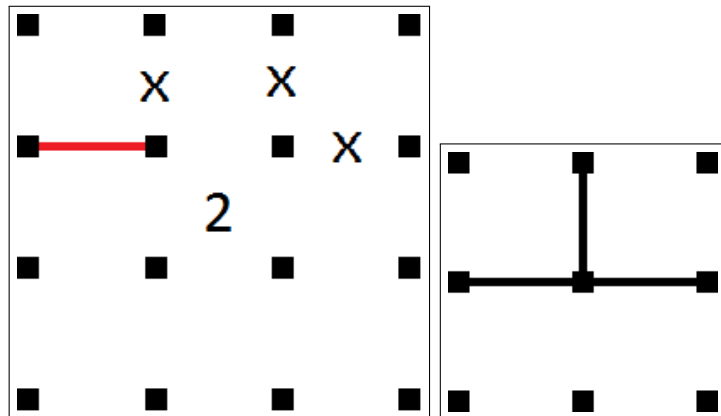


# Static Rules

Static rules are rules that do not contain lines or x's in their prerequisites. We identified 3 static rules.



## Rule and Contradiction in action



We chose to cover rules that are at most 3 by 3 in dimension, and contradictions that are at most 2 by 2 in dimension.

# Contours

- ▶ Use 2D array to keep track of contour endpoints.
- ▶ Update endpoints as we add lines.
- ▶ Keep track of the number of open and closed contours as we add lines.

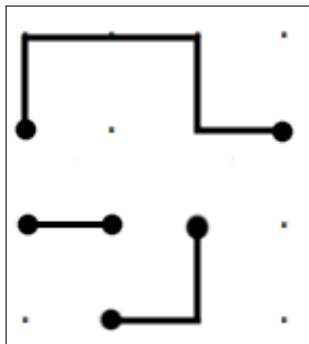


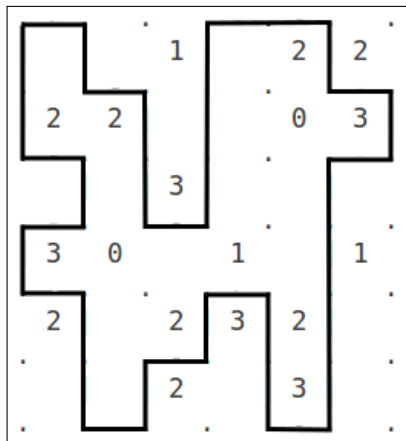
Table 1: Contour Endpoint Array

3,1			0,1
1,2	0,2	1,3	
	2,2		

$numClosed_ = 0$

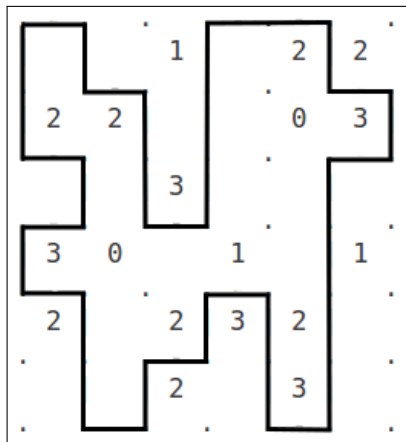
$numOpen_ = 3$

## How Can We Tell Our Grid is Solved?



- ▶ Every number in the grid is satisfied

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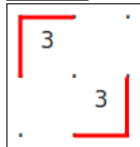
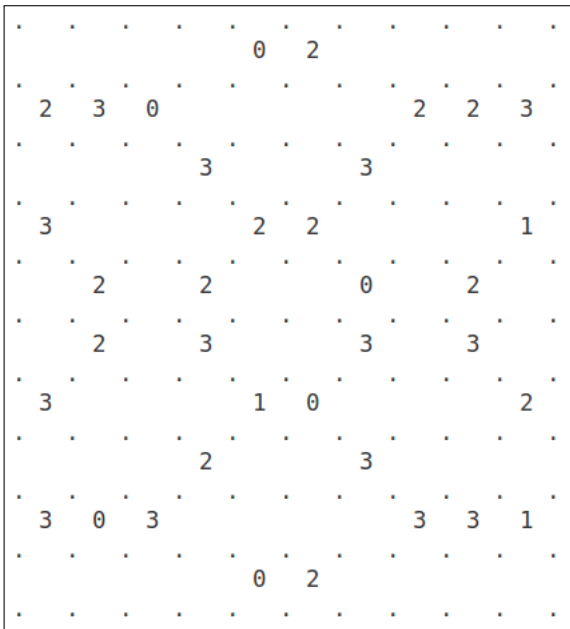
- ▶ Every number in the grid is satisfied
- ▶ There is exactly one closed loop, and no open loops.

# Applying Rules

- ▶ for every position in the grid...
  - ▶ for every defined rule...
    - ▶ for every orientation...

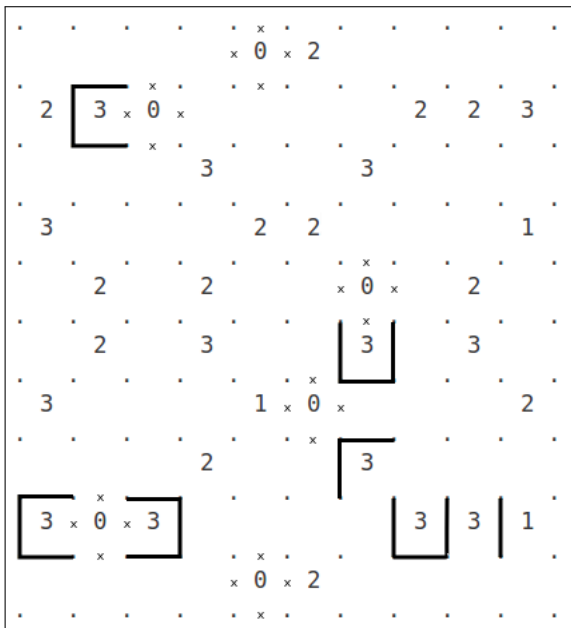
Do the prerequisites in the rule match where we're looking at on the grid?

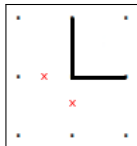
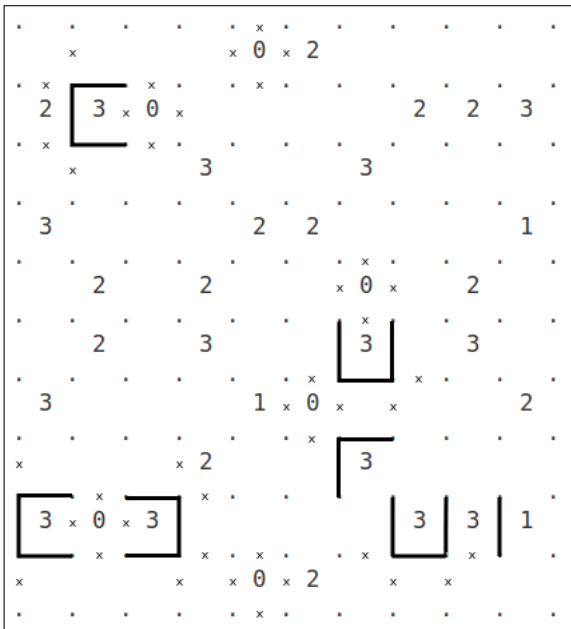
- ▶ If so, add consequences to the grid.

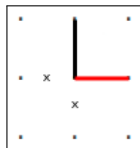
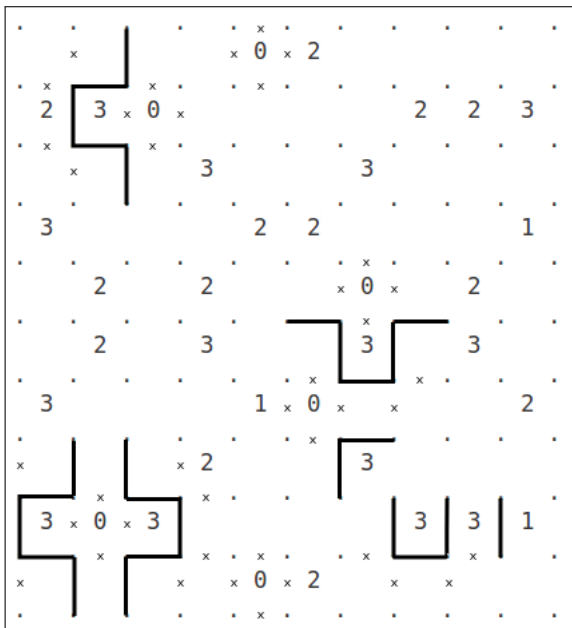


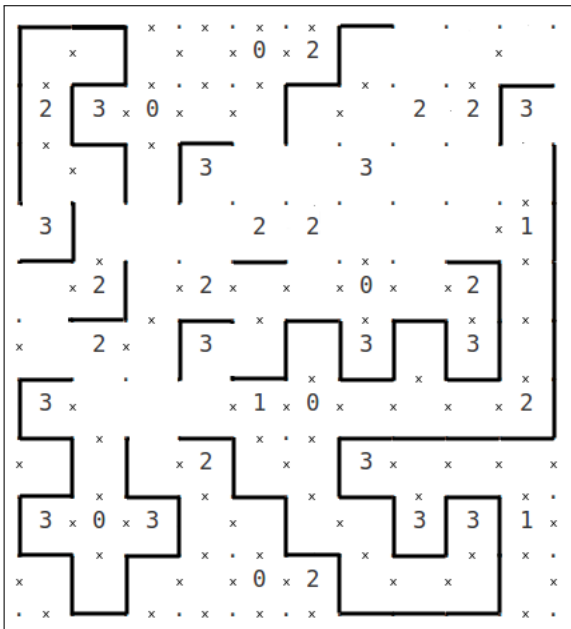












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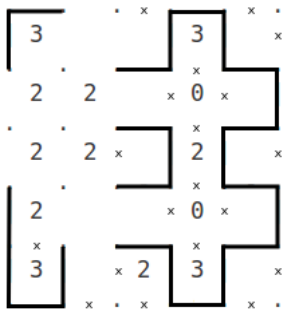
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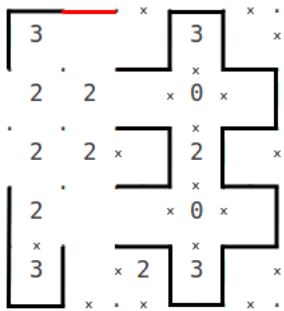
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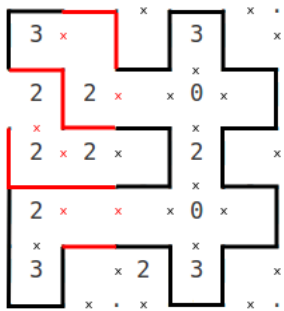
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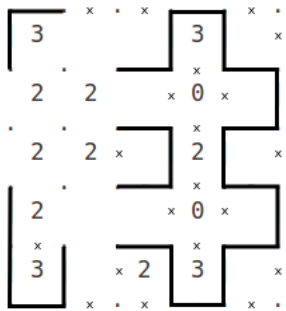
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4. If neither results in a contradiction, take the intersection of the two resulting grids











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If single guesses don't result in anything, we can nest our guesses. New information from nested guesses propagates out to the canonical grid.

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We call  $n$  nested guesses a "depth  $n$  guess".

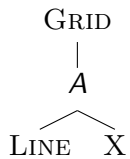
Example of a depth 2 guess

GRID

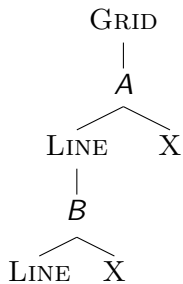
## Example of a depth 2 guess

GRID  
|  
A

## Example of a depth 2 guess

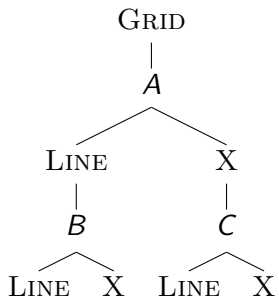


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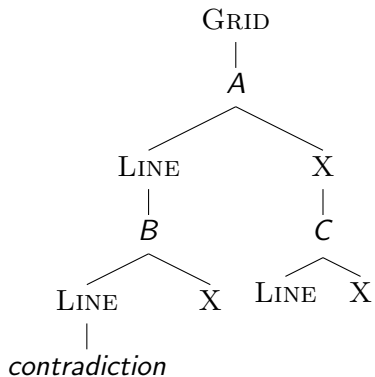




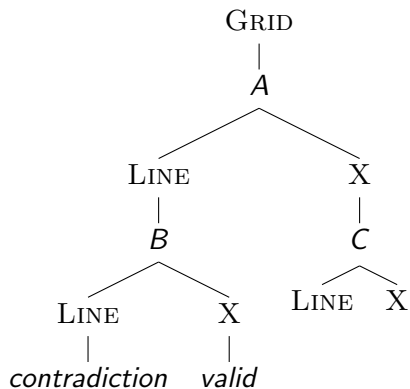
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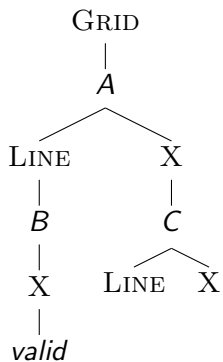
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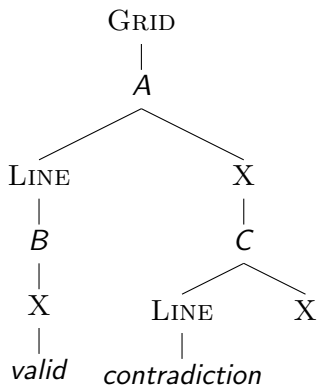
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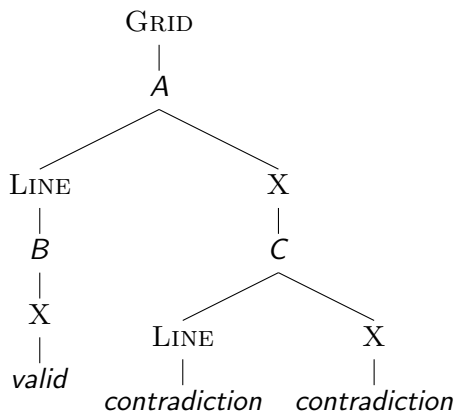
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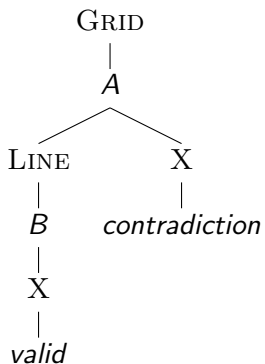
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|

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|

LINE

|

*B*

|

X



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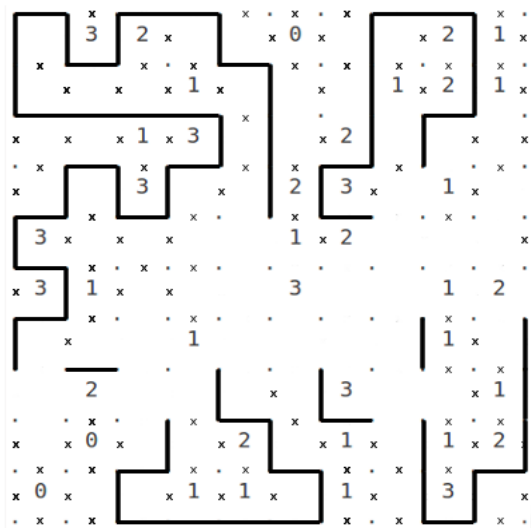
## Repeated guessing algorithm

- ▶ If at any point new information is found, restart algorithm
- ▶ Run deterministic rules
- ▶ Run every possible guess (depth 1 guessing)
- ▶ Run every possible guess, and within each guess, make all possible guesses (depth 2 guessing)

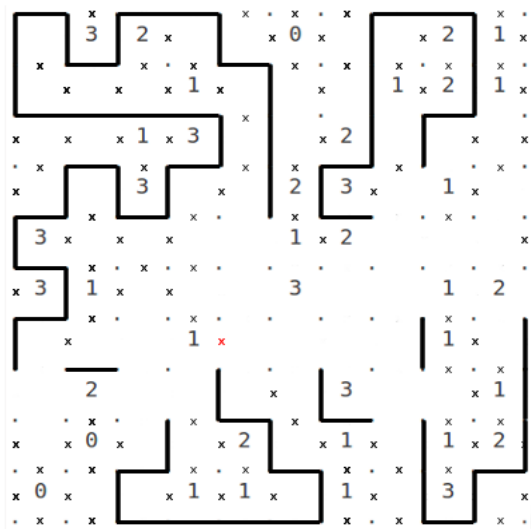
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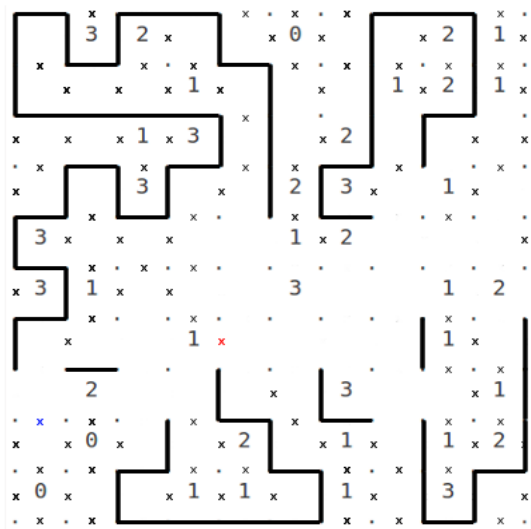
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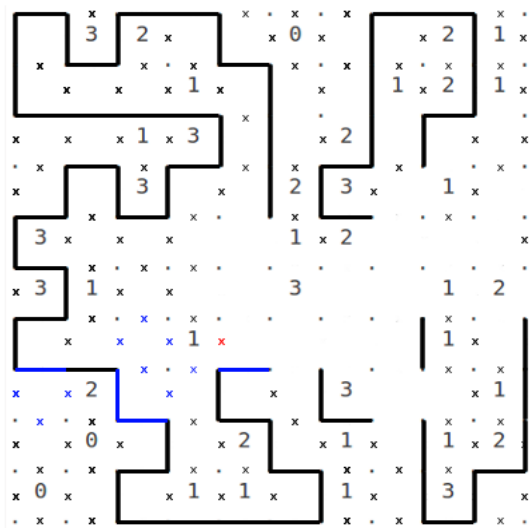


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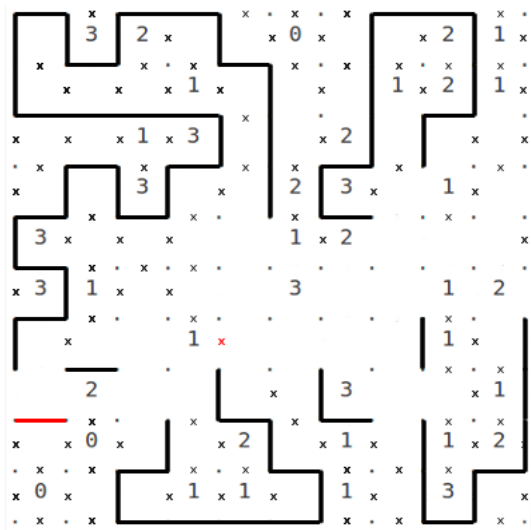




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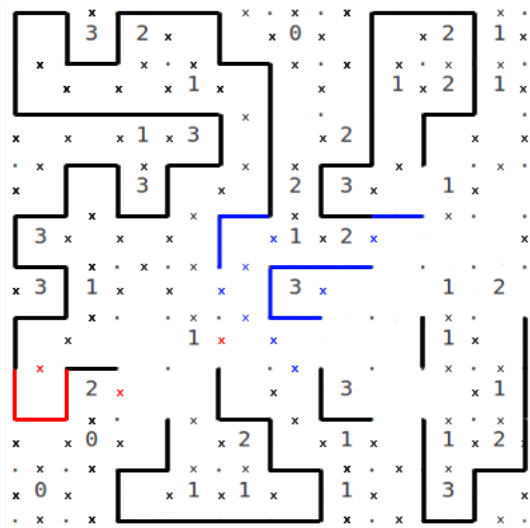
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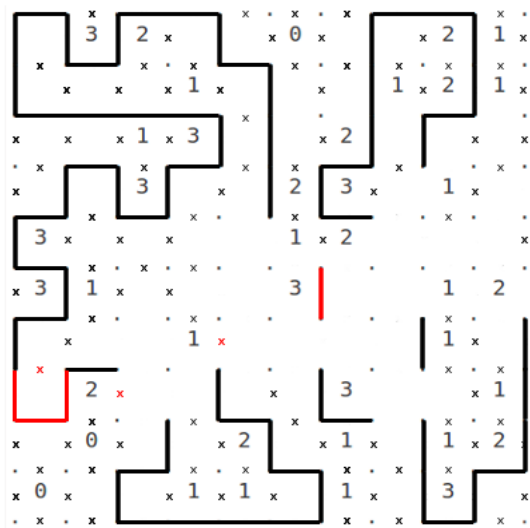




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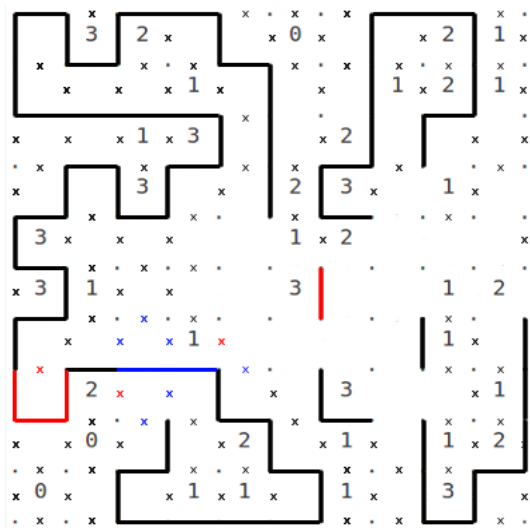


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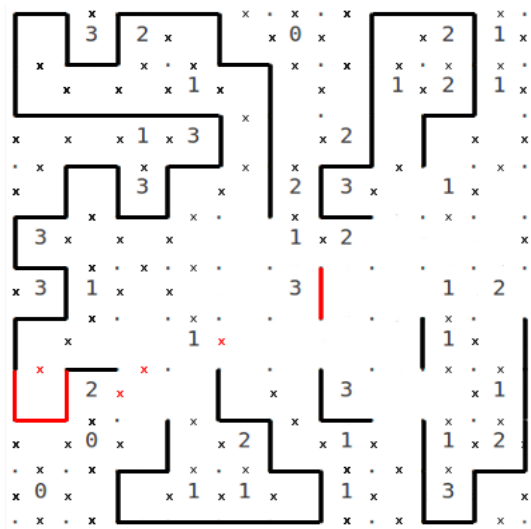


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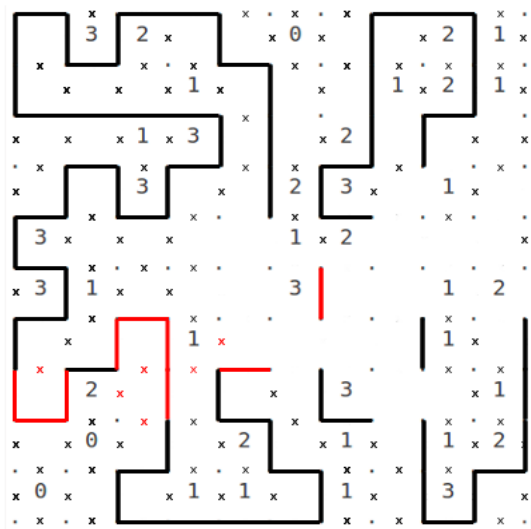




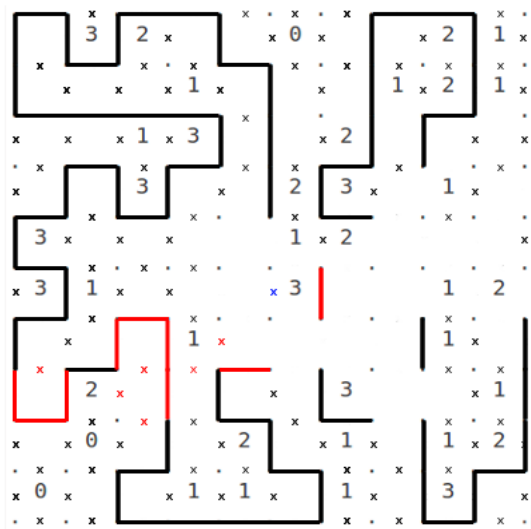
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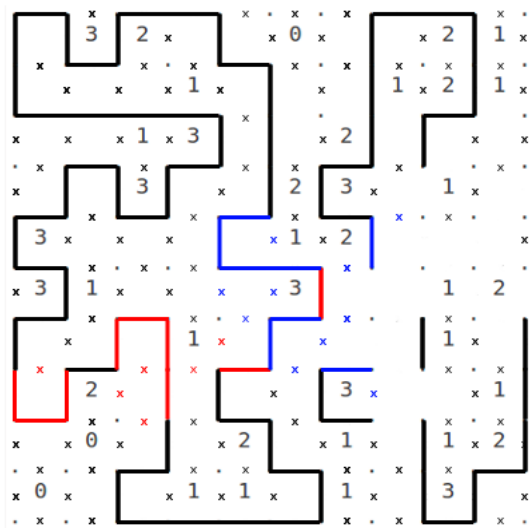
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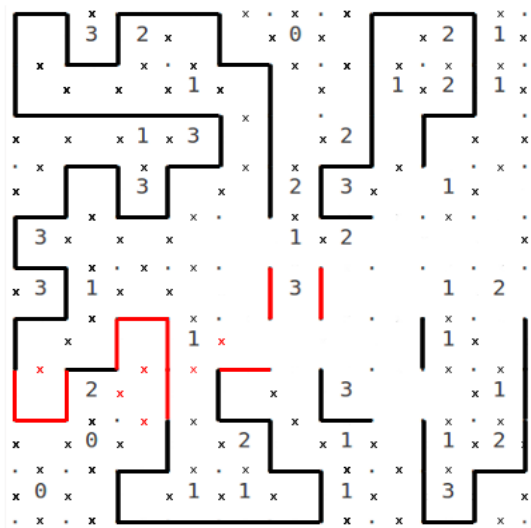
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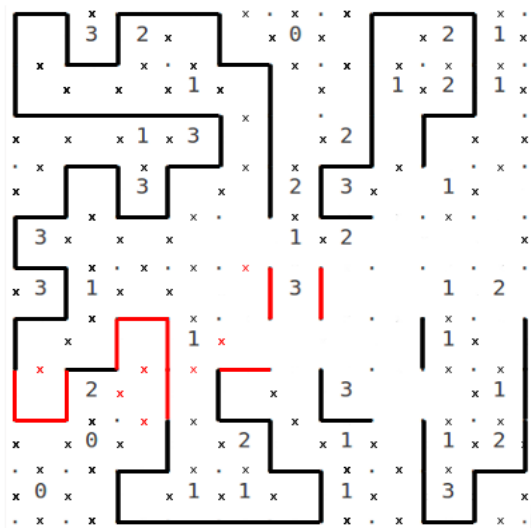
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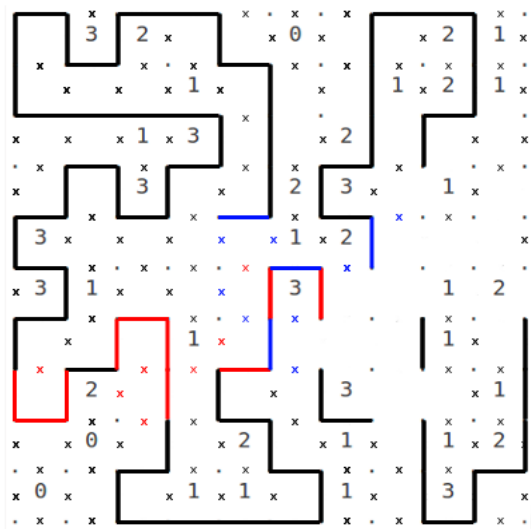


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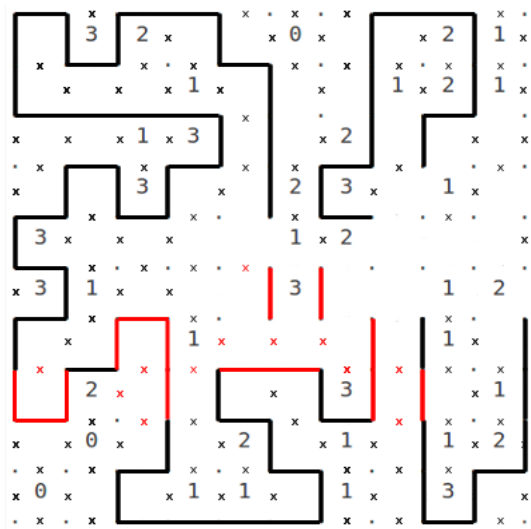
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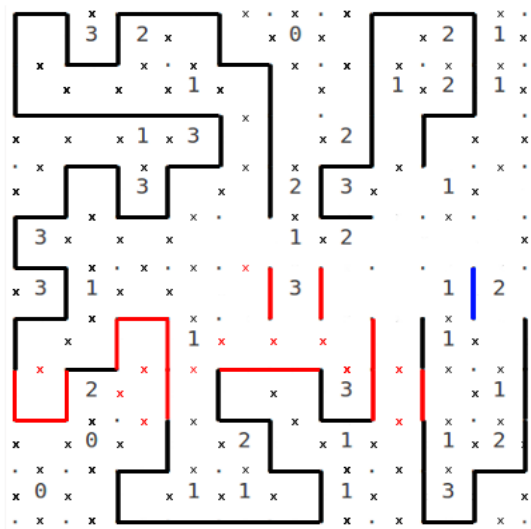




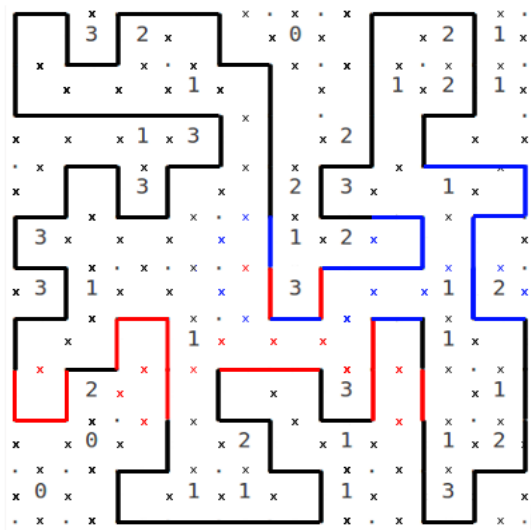
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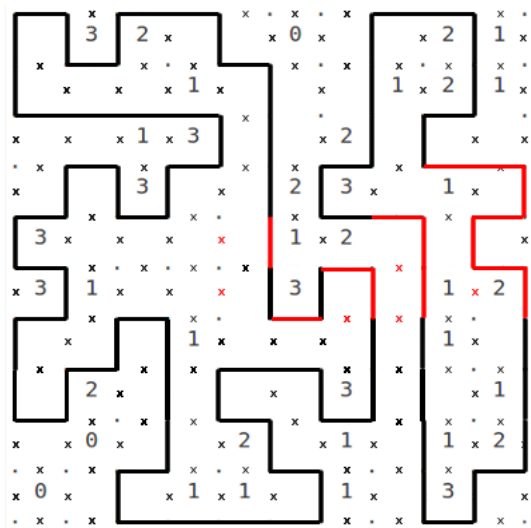
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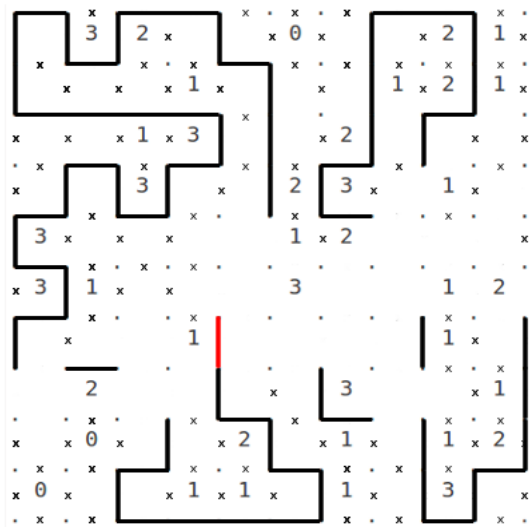
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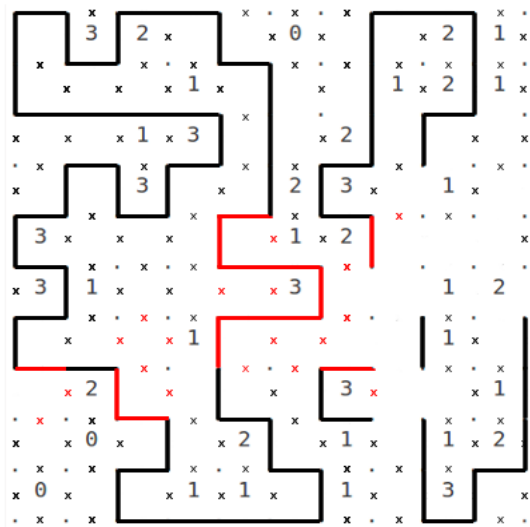
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- ▶ For each depth,  $\mathcal{O}(mn)$  new guesses, each taking  $\mathcal{O}(mn)$  time to instantiate
- ▶ Guessing at the max depth is by far the most important factor in runtime
- ▶ We can maximize this by never filling anything in at lower depths

## Time Complexity

Overall runtime  $\mathcal{O}((mn)^d)$  with  $d$  bounded above by  $\mathcal{O}(mn)$

Table 2: Empty grid completion time

Size	Max Depth	Time (sec)
3x3	0	0.001111
3x3	1	0.033743
3x3	2	1.19998
3x3	3	57.2004
3x3	4	2587.52

## Empirical Results: Typical Puzzles

Table 3: Solve Times

Size	Max Depth	Solve Time
10x10	1	0.048542 seconds
10x10	1	0.385348 seconds
10x10	2	1.77571 seconds
10x10	2	3.23716 seconds
10x10	3	150.824 seconds*
30x25	1	1.95466 seconds
30x25	1	2.38471 seconds
30x25	1	4.4892 seconds
40x30	1	1.97524 seconds
40x30	3	66.268 seconds*

\*These puzzles were determined to have multiple solutions  
Puzzles taken from *nikoli.com*, *kakuro – online.com*, and  
*puzzle – loop.com*.

# Make a Slitherlink Puzzle

## Overview

1. Make a loop
2. Fill grid with numbers
3. Remove some numbers

## Making a Loop

Start with an empty  $m \times n$  grid, a simple rule, and three lists:

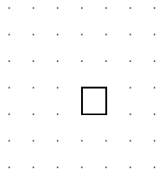
1. *available*
2. *expandable*
3. *unexpandable*

## Making a Loop

Start with an empty  $m \times n$  grid, a simple rule, and three lists:

1. *available*
2. *expandable*
3. *unexpandable*

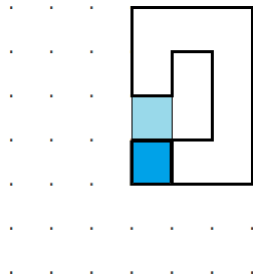
Start with every location in the grid in *available* but none in *unexpandable* and *expandable*. Then, add one random location to *expandable* and remove it from *available*.



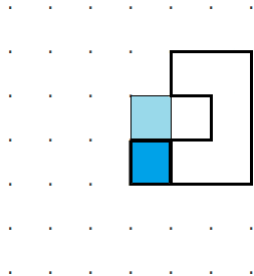
## Bad Stuff

*Rule:* When expanding from a location, *cur*, in *expandable* to an adjacent location in *available*, make sure that adding *pos* to *expandable* doesn't cause any bad stuff

Opposite:



Opposite kitty-corners:



If *opposite* or either of the *opposite kitty-corners* are in *expandable*, then do not add *pos* to the loop.



## Making a Loop cont.

1. Choose a location, *cur* in *expandable* at random

## Making a Loop cont.

1. Choose a location, *cur* in *expandable* at random
2. Look at neighbors to see if and where the loop can expand from *cur*.

## Making a Loop cont.

1. Choose a location, *cur* in *expandable* at random
2. Look at neighbors to see if and where the loop can expand from *cur*.
  - 2.1 If a neighbor is in *available* and it wasn't a valid neighbor, remove it from *available*

## Making a Loop cont.

1. Choose a location, *cur* in *expandable* at random
2. Look at neighbors to see if and where the loop can expand from *cur*.
  - 2.1 If a neighbor is in *available* and it wasn't a valid neighbor, remove it from *available*
  - 2.2 If there are no valid neighbors, add *cur* to *unexpandable* and remove it from *expandable*

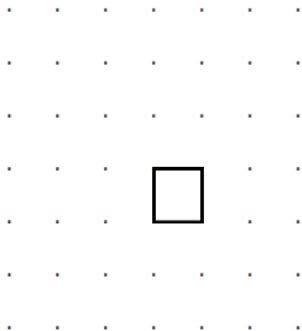
## Making a Loop cont.

1. Choose a location, *cur* in *expandable* at random
2. Look at neighbors to see if and where the loop can expand from *cur*.
  - 2.1 If a neighbor is in *available* and it wasn't a valid neighbor, remove it from *available*
  - 2.2 If there are no valid neighbors, add *cur* to *unexpandable* and remove it from *expandable*
  - 2.3 Otherwise, randomly choose an valid neighbor to add to add to *expandable* and take out of *available*

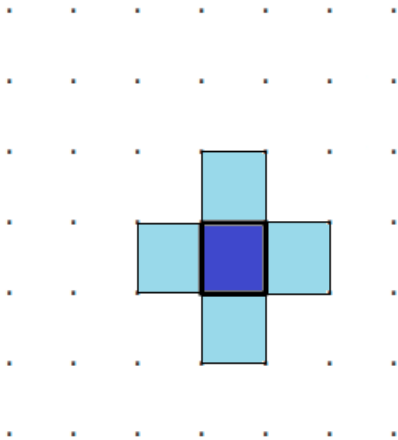
## Making a Loop cont.

1. Choose a location, *cur* in *expandable* at random
2. Look at neighbors to see if and where the loop can expand from *cur*.
  - 2.1 If a neighbor is in *available* and it wasn't a valid neighbor, remove it from *available*
  - 2.2 If there are no valid neighbors, add *cur* to *unexpandable* and remove it from *expandable*
  - 2.3 Otherwise, randomly choose an valid neighbor to add to add to *expandable* and take out of *available*
3. repeat until there are no locations in *expandable*

## Making a Loop Example

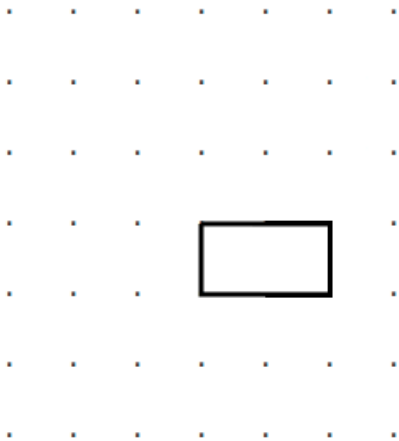


## Making a Loop Example

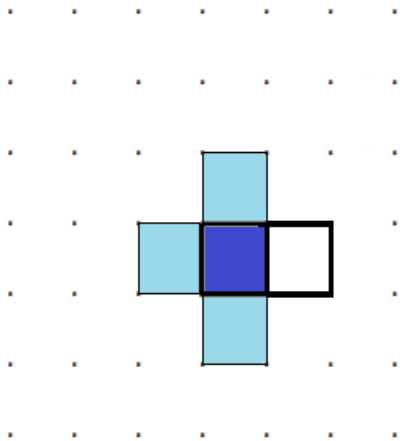




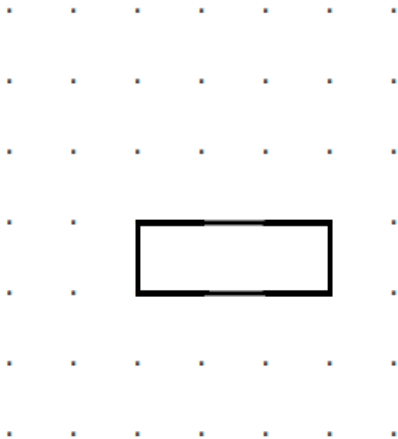
## Making a Loop Example



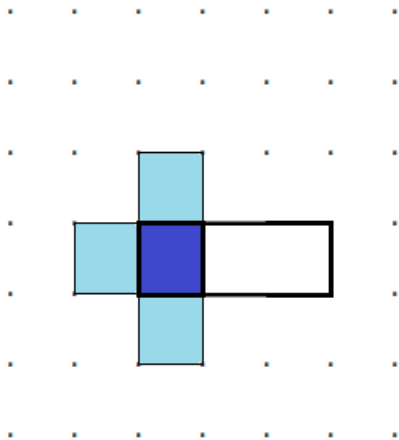
## Making a Loop Example



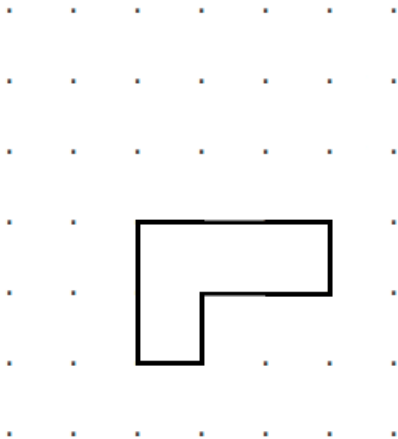
## Making a Loop Example



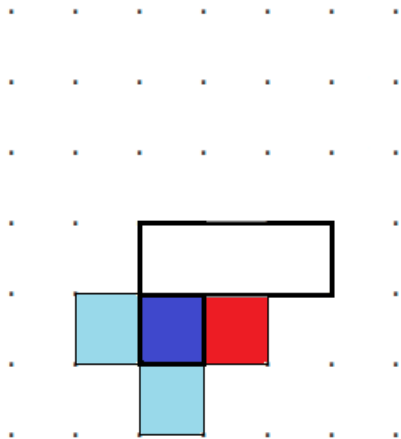
## Making a Loop Example



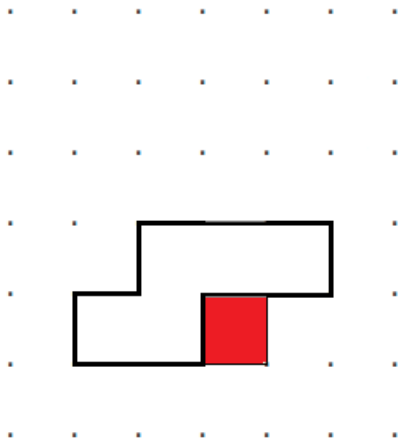
## Making a Loop Example



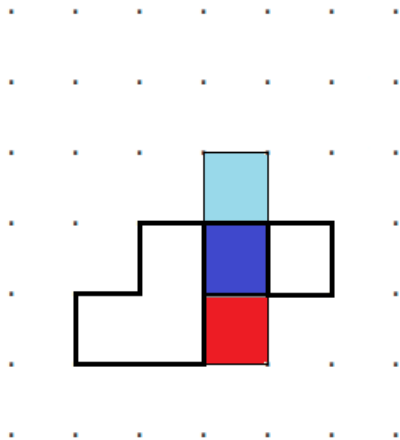
## Making a Loop Example



## Making a Loop Example

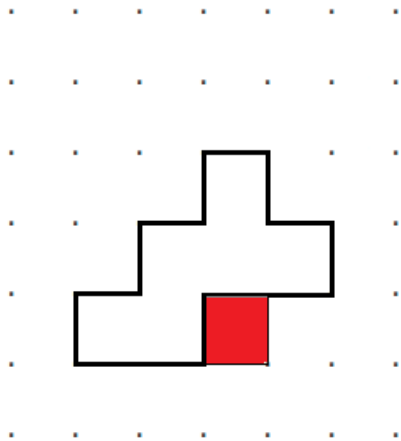


## Making a Loop Example

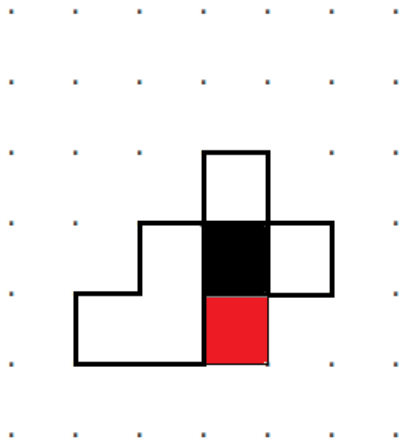




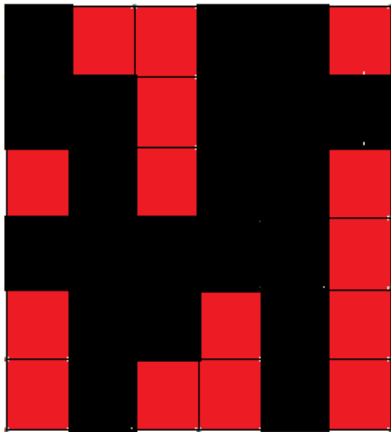
## Making a Loop Example



## Making a Loop Example

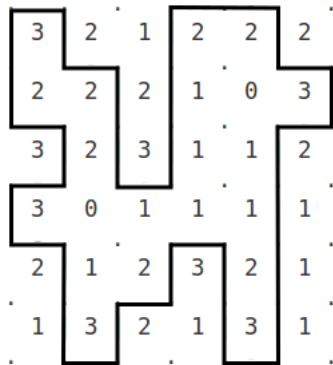
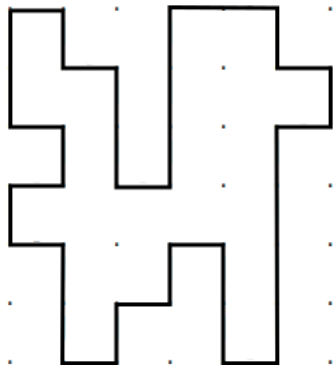


## Making a Loop Example



# Filling with Numbers

Surprise surprise, this is actually really easy



(sum the differences in the neighboring locations)

## Removing Numbers

To make puzzles interesting, we want to remove numbers  
We want to do so until we have reached a certain count  
Must retain one unique solution

# Removing Numbers

The Process:

# Removing Numbers

The Process:

1. Pick a number from a set of eligible numbers

# Removing Numbers

The Process:

1. Pick a number from a set of eligible numbers
2. Add this number to a stack of ineligible numbers



# Removing Numbers

The Process:

1. Pick a number from a set of eligible numbers
2. Add this number to a stack of ineligible numbers
3. Check if eliminating would make the puzzle unsolvable

# Removing Numbers

The Process:

1. Pick a number from a set of eligible numbers
2. Add this number to a stack of ineligible numbers
3. Check if eliminating would make the puzzle unsolvable
  - 3.1 If solvable, remove the number from both the grid and set of eligible numbers

# Removing Numbers

The Process:

1. Pick a number from a set of eligible numbers
2. Add this number to a stack of ineligible numbers
3. Check if eliminating would make the puzzle unsolvable
  - 3.1 If solvable, remove the number from both the grid and set of eligible numbers
  - 3.2 If unsolvable, only remove the number from the set of eligible numbers

# Removing Numbers

The Process:

1. Pick a number from a set of eligible numbers
2. Add this number to a stack of ineligible numbers
3. Check if eliminating would make the puzzle unsolvable
  - 3.1 If solvable, remove the number from both the grid and set of eligible numbers
  - 3.2 If unsolvable, only remove the number from the set of eligible numbers
4. Repeat until set of eligible numbers is empty

## Removing Numbers cont.

Once set of eligible numbers is empty:

## Removing Numbers cont.

Once set of eligible numbers is empty:

1. Pop numbers off ineligible stack

## Removing Numbers cont.

Once set of eligible numbers is empty:

1. Pop numbers off ineligible stack
2. Place each back into the set of eligible numbers

## Removing Numbers cont.

Once set of eligible numbers is empty:

1. Pop numbers off ineligible stack
2. Place each back into the set of eligible numbers
3. Do so until most recently eliminated number is found



## Removing Numbers cont.

Once set of eligible numbers is empty:

1. Pop numbers off ineligible stack
2. Place each back into the set of eligible numbers
3. Do so until most recently eliminated number is found
4. Keep eliminated in the ineligible stack, but place back into grid

## Removing Numbers cont.

Once set of eligible numbers is empty:

1. Pop numbers off ineligible stack
2. Place each back into the set of eligible numbers
3. Do so until most recently eliminated number is found
4. Keep eliminated in the ineligible stack, but place back into grid

Repeat removing numbers until desired count is reached

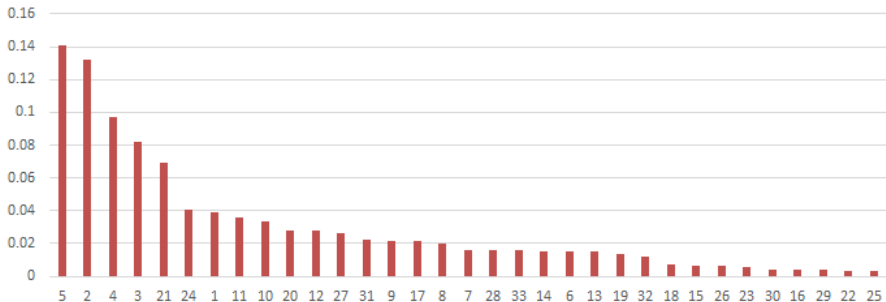
## Removing Numbers cont.

It's too hard!

### Improvements

1. Data
2. Rule set
3. Balancing

## Rule Frequency, sorted



## Ruleset Limitation

As a result, we created two subsets of rules:

- ▶ easy: the rules with a greater than 5% occurrence rate
- ▶ hard: the rules with a greater than 1% occurrence rate

## An 'Easy' Puzzle

.	.	.	.	.	.	.	.
			2	1	2	2	
.	.	.	.	.	.	.	.
	2	0		3	1	1	
.	.	.	.	.	.	.	.
				1			
.	.	.	.	.	.	.	.
	1	2	3	1		1	
.	.	.	.	.	.	.	.
				1		2	
.	.	.	.	.	.	.	.
	1	1	3	1			
.	.	.	.	.	.	.	.



# Balancing the Numbers

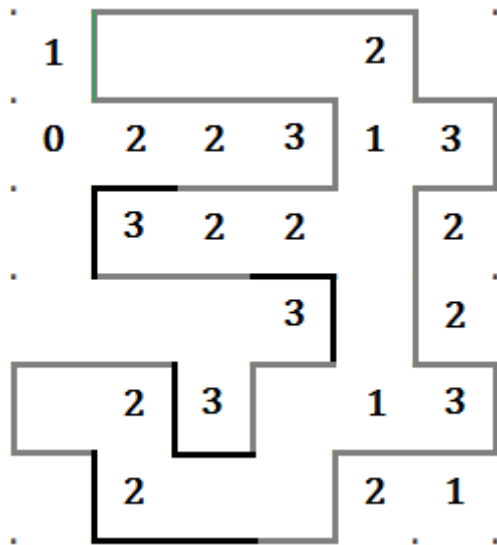




## A Balanced Easy Puzzle

.	.	.	.	.	.	.	.
	<b>1</b>				<b>2</b>		
.	.	.	.	.	.	.	.
	<b>0</b>	<b>2</b>	<b>2</b>	<b>3</b>	<b>1</b>	<b>3</b>	
.	.	.	.	.	.	.	.
		<b>3</b>	<b>2</b>	<b>2</b>		<b>2</b>	
.	.	.	.	.	.	.	.
				<b>3</b>		<b>2</b>	
.	.	.	.	.	.	.	.
		<b>2</b>	<b>3</b>		<b>1</b>	<b>3</b>	
.	.	.	.	.	.	.	.
		<b>2</b>			<b>2</b>	<b>1</b>	
.	.	.	.	.	.	.	.

## A Quick Attempt at the 'Easy' Puzzle



# Time Complexity

## Important details

- ▶ The Solver is run on the order of  $mn$  times.
- ▶ Each time the solver is run, it happens with a maximum depth of one guess which has on the order of  $\mathcal{O}((mn)^2)$  time.
- ▶ Therefore, the generator run in the order of  $\mathcal{O}((mn)^3)$  time.

## References

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*On the NP-completeness of the Slither Link Puzzle* Takayuki YATO  
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*A rule-based approach to the puzzle of Slither Link.* Stefan Herting.

*Puzzles and Games: A Mathematical Modeling Approach.* Tony Hürlimann, 2015

*Solving logical puzzles using mathematical models.* KVIS Susanti, S Lukas.

Thank you

Questions?